

## **ANONYMITY IN** QUANTUM NETWORK PROTOCOLS





## Fit into QIA's bigger picture

- **Application domains**, User domains and use-cases ullet
- **Identity protection** •









A quantum network  $N = \{1, 2, ..., n\}$ 

- The 'participants':
  - A sender: Alice A
  - Multiple *receivers:* Bobs  $B_i$
- The 'non-participants'
  - Everyone else











The identity of participants remains "as unknown as it was"

- To the non-participants: weak anonymity ullet
- To everyone\*: strong anonymity •
  - Alice chooses all identities
  - Variants  $\bullet$







ANONYMOUS SECRET MESSAGING

### Anonymous secret messaging

- **1.** Alice picks the Bobs
- **2.** Alice and the Bobs anonymously generate key k
- **3.** Alice encrypts her message  $s \rightarrow c = s \oplus k$
- **4.** Alice anonymously broadcasts c
- **5.** Bobs decrypt  $c \to c \oplus k = m \oplus k \oplus k = m$









#### **GHZ** states

## $|0...0\rangle_N + |1...1\rangle_N$

- Measure in Z•
  - Correlated key •
- Measure in  $X_n$ 
  - 'GHZ' on remaining parties
- Measure **all qubits** in X
  - The *parity* of all outcomes is fixed







#### $(|0...0\rangle_{N-1} \pm |1...1\rangle_{N-1}) \otimes |\pm\rangle_n$

(GHZ is stabilizer)



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## BROADCASTING



Protocol for Alice to send message b

- **1.** Share GHZ state over entire network  ${\cal N}$
- **2.** Alice applies  $Z_a$  if b = 1
- **3.** Everyone measures  $X_i$ , obtain  $m_i$ , announce
- **4.** Now  $b = \bigoplus m_i$





## $|0...0\rangle + (-1)^{0}|1...1\rangle$ $|0...0\rangle + (-1)^{b_{a}}|1...1\rangle$



## **KEY AGREEMENT**



#### What we want









#### Verification

- Bobs select random X and Y measurements
- Alice completes a stabilizer element
- **Parity** is fixed

• **Repeating** gives exponentially small cheat probability







#### **Contribution to QIA and future steps**

Application domain

• Crystallise a Use case instead

Invitation to tomorrow morning







#### Links

#### Anonymous Tranmissions: lacksquare

- Christandl, Wehner (2005). Quantum Anonymous Transmissions. <u>https://doi.org/10.1007/11593447\_12</u> •
- Anonymous conference key agreement: ullet
  - Hahn, de Jong, Pappa (2020). Anonymous Quantum Conference Key Agreement. https://doi.org/10.1103/PRXQuantum.1.020325 ۲
  - Grasselli, Murta et al. (2022). Secure anonymous conferencing in quantum networks. https://doi.org/10.1103/PRXQuantum.3.040306 •
  - de Jong, Hahn et al. (2020). Anonymous conference key agreement in linear quantum networks. https://arxiv.org/abs/2205.09169 •

#### Experimental work: $\bullet$

- Thalacker et al. (2021). Anonymous and secret communication in quantum networks. https://doi.org/10.1088/1367-2630/ac1808 ۲
- Rückle et al. (2022). Experimental anonymous conference key agreement using linear cluster states. https://arxiv.org/abs/2207.09487 ٠







# THANK YOU!

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(Bad) Protocol for Alice and Bobs to agree on key k

- **1.** Share GHZ state over entire network  $\boldsymbol{N}$
- **2.** All non-participants j measure  $X_j$
- **3.** All participants *j* measure  $Z_j$ , obtain  $k_j$
- **4.** All participants agree on  $k = k_i$

What happens when non-participants deviate?





# $|0...0\rangle_N + (-1)^0 |1...1\rangle_N$ $|0...0\rangle_P + (-1)^\alpha |1...1\rangle_P$



Protocol for Alice and Bobs to agree on key  $\boldsymbol{k}$ 

- **1.** Share GHZ state over entire network  ${\cal N}$
- **2.** All non-participants j measure  $X_j$
- 3. Network asks public random source for random  $\boldsymbol{b}$
- 4a. (Key generation)

Participants measure Z and obtain key k

4b. (Verification)

Participants run verification round





# Repeat ad inf.



(Supposedly) the state is:  $(|0...0\rangle_P + (-1)^{\alpha} |1...1\rangle_P) \otimes |garb\rangle_{NP}$ 

- **1.** All Bobs pick random  $b_i$ , measure  $X_i$  or  $Y_i$
- **2a.** All Bobs announce  $b_i$  and outcome  $o_i$
- **2b.** Others announce random  $b_i$  and  $o_i$
- **3.** Alice resets  $b_a = \bigoplus_{Bobs} b_j$ , measures  $X_j$  or  $Y_j$
- **4.** Alice accepts iff  $o_a \oplus \alpha = \bigoplus_{Bobs} o_i$







#### Important points

- **1.** Timing issues
- 2. Initial GHZ state is not checked
- 3. Dephasing of GHZ results in failed verification





